RFC: Direct Chunk Write

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This document discusses a new HDF5 C function to bypass the library’s data conversion and filter pipeline and write data chunks directly to a dataset in the file. It recommends putting the new function in the HDF5 high-level C library.

1 Introduction

A customer of the HDF Group requested a new function, which would allow an application to write pre-compressed data chunks directly to a dataset in an HDF5 file. The application will compress its data outside the HDF5 library so that it is no longer limited by HDF5 library’s filter pipeline.

2 Motivation

The customer designs next generation X-ray pixel detectors at synchrotron light sources. It produces data at the rate of tens of gigabyte per second. The modular architecture of the detector can scale up its data stream in parallel. It maps well to the current parallel computing and storage systems. The data stream can be compressed well. The data volume can be reduced by a factor of ten or more after compression. Hence, it is crucial to compress the data before storing on disk. Any file format for storing detector data must support data compression and allow the scaling up of the number of processing nodes.

The current HDF5 library restricts compression to its filter pipeline and thus limits the performance within a single process. To overcome this weakness, the library should provide a new function to allow applications to write compressed data chunks directly to a dataset in an HDF5 file. It allows the application to compress the detector data in parallel and match the data production rate of the detector.
3 Requirements

The customer gave us the following requirements for the function.

3.1 Functionality

The new function should

- Accept converted and compressed data chunks in a memory buffer
- Provide a parameter for a chunk index that identifies the chunk's position in its dataset
- Provide a parameter to indicate which filters of the dataset's filter pipeline should be skipped when reading the chunk back from the file
- Accept a simple value 0 for the filter status to indicate that no filter shall be skipped in order to simplify the application
- Be available in a release of version 1.8 and 1.10

3.2 Implementation

The implementation of the new function in the library should

- Use the application’s original data buffer in memory to write the data chunk in the file without copying the data to any internal buffer
- Apply no modification or check to the data in the application’s memory buffer
- Accept from the application data of any arbitrary non-zero size
- Allow the application to reuse the memory buffer once the library returns the function
- Allow the writing of a partial chunk on the edge of a dataset
- Allow the overwriting of existing data chunks
4 Possible Solutions

The new function differs from the current HDF5 H5D interface functions. It takes applications’ data and writes it directly to a file bypassing the HDF5 library’s internals. This is the first function of this type that The HDF Group developers were tasked to implement. Current library architecture doesn’t have any good place for adding such functions. Changes to the HDF5 library architecture are required.

This section discusses three possible solutions for the new function: in the HDF5 C library H5D interface, in the HDF5 high-level C library, or in the HDF5 C library but in a new library source file.

4.1 In the HDF5 C Library

The first solution proposes that the new function is in the library among other HDF5 H5D C API functions. Below is the function prototype (do stands for dataset optimization):

```c
herr_t H5DOwrite_chunk(hid_t dset_id,
                         hid_t dxpl_id,
                         uint32_t filter_mask,
                         hsize_t *offset,
                         size_t data_size,
                         const void *buf)
```

Please see Appendix B for the full description of this function. Basically, this function takes a pre-processed data chunk (buf) and its size (data_size) and writes to the chunk location (offset) in the dataset (dset_id).

A simple example of using H5DOwrite_chunk to write one chunk of data shown on Figure 1 is like the following (a more complete example is in Appendix B):

```c
hsize_t offset[2] = {4, 4};
uint32_t filter_mask = 0;
size_t nbytes = 40;
if(H5DOwrite_chunk(dset_id, dxpl, filter_mask, offset, nbytes, outbuf) < 0)
    goto error;
```

In this example, the dataset is 8x8 elements of int. Each chunk is 4x4. The offset of the first element of the chunk to be written is 4 and 4. In the diagram below (Figure 1), the shaded chunk is the data to be written.
Figure 1. Illustration of the chunk to be written in the example code above

![Diagram](image)

The function is writing a pre-compressed data chunk of 40 bytes (assumed) to the dataset. The offset of first element of the chunk is 4 and 4. The zero value of the filter mask means that all filters have been applied to the pre-processed data.

**Discussion**

This approach was implemented when we tested the feasibility of the function and performance. The benchmarks ran by the customer and The HDF Group developers showed order of magnitude enhancement in I/O. Implementation met the requirements and expectation of the customer.

Integrating this function into the library does not require much work. Using this function will evoke little overhead as it deals with the low-level routines inside the library.

However, to use it appropriately, one has to understand what this function exactly does – bypassing data conversion, chunk cache, filter pipeline, etc. Since it touches the low-level library directly, it lacks the protection that most public functions have. It is more specialized compared to other general public functions. It requires the user to have special knowledge to use this function properly.

Maintenance of this function will require some careful considerations since it has a different architecture and may be confusing to the future maintainers of the HDF5 H5D interface.

**Design**

Figure 2 below shows that using `H5Dwrite_chunk` to write pre-compressed data chunks can bypass the library's data conversion and filter pipeline.

**Implementation**

In the internal routine `H5D_chunk_direct_write` of `H5Dchunk.c` under the library’s source directory, the functionality of direct chunk write is achieved in the following steps:

1. Initialize the chunked storage (such as B-tree indexing for chunks) through `H5D__alloc_storage` if it has not been done yet.
2. Calculate the chunk index with the chunk offset passed in as a parameter through `H5V_chunk_index`. 

3. Find out the file address of the chunk to be written through \texttt{H5D\_chunk\_lookup}.

4. Create the chunk if it does not exist, or re-allocate the chunk if its size changes through a callback function for insertion.

5. Evict the entry for the chunk through \texttt{H5D\_chunk\_cache\_evict} if it is in the chunk cache without flushing it to the file.

6. Write the chunk to the file through \texttt{H5F\_block\_write}.
Performance

The following table (Table 1) is the result of the performance benchmark test run by the HDF developers. It shows that using the new function `H5DOWrite_chunk` to write pre-compressed data is much faster than using the `H5Dwrite` function with the filter pipeline. The test was run on a Linux 2.6.18 / 64-bit intel x86_64 machine (koala). In the table below we highlighted results that show performance improvement for writing a compressed chunk using `H5DOWrite_chunk`. 
Table 1. Performance result for H5DOwrite_chunk

<table>
<thead>
<tr>
<th>Dataset size (MB)</th>
<th>95.37</th>
<th>762.94</th>
<th>2288.82</th>
</tr>
</thead>
<tbody>
<tr>
<td>Data dimension (number of elements)</td>
<td>100x1000x250</td>
<td>100x2000x1000</td>
<td>100x2000x3000</td>
</tr>
<tr>
<td>Chunk dimension (number of elements)</td>
<td>1000x250</td>
<td>2000x1000</td>
<td>2000x3000</td>
</tr>
<tr>
<td>H5DOwrite_chunk with uncompressed data (MB/s)</td>
<td>76.99</td>
<td>96.24</td>
<td>94.72</td>
</tr>
<tr>
<td>H5DOwrite_chunk with pre-compressed data (MB/s)</td>
<td>77.31</td>
<td>94.71</td>
<td>96.91</td>
</tr>
<tr>
<td>H5Dwrite without compression filter (MB/s)</td>
<td>76.22</td>
<td>95.66</td>
<td>95.4</td>
</tr>
<tr>
<td>H5Dwrite with compression filter (MB/s)</td>
<td>2.68</td>
<td>2.68</td>
<td>2.66</td>
</tr>
<tr>
<td>Write pre-compressed data to a Unix file (MB/s)</td>
<td>78.25</td>
<td>95.75</td>
<td>98.16</td>
</tr>
</tbody>
</table>

*Each element is a 4-byte integer

4.2 In the HDF5 High-level C Library

The second solution proposes that the new function will be put into the high-level library. The function is still called H5DOwrite_chunk. It is used in the same way as the first proposal (4.1).

Discussion

This approach requires additional considerations to ones discussed in (4.1). Following the convention that all high-level functions use only the library’s public functions, this function should call H5Dwrite. In order to call H5Dwrite the following additional steps are taken: the parameters for direct chunk write function H5DOwrite_chunk will be saved in the dataset transfer property list and be transferred into H5D_pre_write (a new internal function). In H5D_pre_write, the library will retrieve the information for direct chunk write from the dataset transfer property list and handle it as a special case. This approach requires the re-evaluation of the performance achieved with 4.1, especially from the customer’s side.

Design

Figure 3 below shows the design layout for H5DOwrite_chunk in the high-level library and the subsequent changes of internal functions in the library source code.

The functionality of H5D_chunk_direct_write has been discussed in section 4.1. It bypasses the data gathering and scattering, data conversion, filter pipeline, and chunk cache and writes the data chunk to the file directly. This diagram omits this part of design. Please see Figure 1 in section 4.1.2 for the detail.
Implementation

The functionality of `H5Dwrite_chunk` is implemented in the following steps:

1. The `H5Dwrite_chunk` function in the high-level library sets a flag in the dataset transfer property list to indicate direct chunk write. It also saves its parameters `filter_mask`, `offset`, and `data_size` in the dataset transfer property list. Then it calls `H5Dwrite` since high-level functions can only call the library’s API function.

2. Both regular writes and direct chunk writes go through `H5Dwrite`. `H5D__pre_write` (in `H5Dio.c`) is a new internal function that prepares for calling `H5D__write` or `H5D__chunk_direct_write`.

3. When `H5D__pre_write` sees the flag for direct chunk write is on, it retrieves the information (`filter_mask`, `offset`, and `data_size`) for direct chunk write from the dataset transfer property list and invokes `H5D__chunk_direct_write` (in `H5Dchunk.c`).

4. If it is a regular write, the direct chunk write flag is off. `H5D__pre_write` would then invoke `H5D__write` for a regular write.
In the figure above, DXPL is short for the dataset transfer property list.
Performance

The performance of H5DOWrite_chunk in the high-level library is similar to the function’s performance in the HDF5 library (see 4.1.4). The customer has confirmed that its performance is as good as they expect. In the table below we highlighted results that show performance improvement for writing a compressed chunk using H5DOWrite_chunk. The results are consistent with the results provided in Table 1.

Table 2. Performance result for H5DOWrite_chunk in the high-level library

<table>
<thead>
<tr>
<th>Dataset size (MB)</th>
<th>95.37</th>
<th>762.94</th>
<th>2288.82</th>
</tr>
</thead>
<tbody>
<tr>
<td>Data dimension (number of elements*')</td>
<td>100x1000x250</td>
<td>100x2000x1000</td>
<td>100x2000x3000</td>
</tr>
<tr>
<td>Chunk dimension (number of elements*')</td>
<td>1000x250</td>
<td>2000x1000</td>
<td>2000x3000</td>
</tr>
<tr>
<td>H5DOWrite_chunk with uncompressed data (MB/s)</td>
<td>79</td>
<td>95.71</td>
<td>89.17</td>
</tr>
<tr>
<td>H5DOWrite_chunk with pre-compressed data (MB/s)</td>
<td>77.19</td>
<td>78.56**</td>
<td>96.28</td>
</tr>
<tr>
<td>H5Dwrite without compression filter (MB/s)</td>
<td>77.27</td>
<td>97.02</td>
<td>91.77</td>
</tr>
<tr>
<td>H5Dwrite with compression filter (MB/s)</td>
<td>2.68</td>
<td>2.67</td>
<td>2.67</td>
</tr>
<tr>
<td>Write pre-compressed data to a Unix file (MB/s)</td>
<td>76.49</td>
<td>95</td>
<td>98.59</td>
</tr>
</tbody>
</table>

* Each element is a 4-byte integer

** This IO speed needs further investigation since it is not consistent with the similar result in Table 1.

4.3 In the HDF5 C Library but in a New Library Source File

The third solution proposes that the new function would behave in the same way as the first proposal, but that the function would be located in a separate source file.

The new source file called H5Dopt.c. In the future, other similar optimization functions will be located in H5<X>opt.c. Users can choose to compile this file with the function in it through configuration. This approach needs less work from the first round of prototype testing. The implementation is simpler than the second approach (4.2).

From the user’s perspective, this approach is the same as the first approach. Users need to understand the function’s behaviors – how it bypasses data conversion, chunk cache, and file pipeline.
5 Recommendation

Since the performance of the “direct chunk write” feature is not affected by where the function is located, we recommend that the second proposal be adopted. The extra protections afforded by the high-level library make the extra work worthwhile. This approach is also applicable to other functions that will be added to the HDF5 software in the nearest future thus providing uniform approach for the functions that do not follow the mainstream architecture.
## Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Description</th>
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<tbody>
<tr>
<td>November 14, 2012</td>
<td>Version 1 circulated for comment within The HDF Group.</td>
</tr>
<tr>
<td>November 28, 2012</td>
<td>Version 2 circulated for comment within The HDF Group.</td>
</tr>
<tr>
<td>November 30, 2012</td>
<td>Version 3 sent to the customer</td>
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</table>
Appendix A: Background Material

In June 2012, the user provided us a requirement document called Extension to the HDF5 API and Library: Write Pre-compressed Chunk Data.
Appendix B: Reference Manual Function Entry for H5DOwrite_chunk

Last modified: 14 November 2012

Name: H5DOwrite_chunk

Signature:

```
  herr_t H5DOwrite_chunk( hid_t dset_id, hid_t dxpl_id, uint32_t filter_mask, hsize_t *
  offset, size_t data_size, const void * buf )
```

Purpose:

Writes a raw data chunk from a buffer directly to a dataset.

Description:

H5DOwrite_chunk writes a raw data chunk from the application memory buffer buf directly to its logical destination offset in a dataset identified by dset_id. Typically, the data in buf is preprocessed in memory by a custom transformation, such as compression. The data chunk will bypass the library’s internal data transfer pipeline (including filters), and will be written to the file directly.

dxpl_id is the identifier of a data transfer property list. (Currently, it is unused.)

filter_mask is a mask to keep a record of which filters are used and is saved with the data chunk in the file. The default value is zero and indicates that all enabled filters are applied. A filter is skipped if the bit corresponding to the filter position (0<= * <32) in the pipeline is turned on.

offset is an array that represents the logical position of the first element of the data chunk in the dataset’s dataspace. The length of the offset array must equal the rank (number of dimensions) of the dataspace. The values in offset must not exceed the dimension limits and must fall on the boundary of data chunks.

data_size is the size of the raw data chunk and represents the number of bytes to be read from the buffer buf. If the data chunk has been pre-compressed, it should be the size of the compressed data.

buf is the memory buffer for the data chunk.

Parameters:

- **hid_t dset_id** IN: Identifier of the dataset to write to.
- **hid_t dxpl_id** IN: Identifier of a transfer property list for this I/O operation.
- **uint32_t filter_mask** IN: Mask for filters.
- **hsize_t * offset** IN: Logical position of the chunk’s first element in the dataspace.
- **size_t data_size** IN: Size of the actual data.
- **const void * buf** IN: Buffer with data to be written to the file.
Returns:

Returns a non-negative value if successful; otherwise returns a negative value.

Example Usage:

An example of using `H5DOWrite_chunk` to write an entire dataset by chunk is like the followings:

```c
#include <zlib.h>
#include <math.h>
#define DEFLATE_SIZE_ADJUST(s) (ceil(((double)(s))*1.001)+12)

size_t buf_size = CHUNK_NX*CHUNK_NY*sizeof(int);
const Bytef *z_src = (const Bytef*)(direct_buf);
Bytef *z_dst; /*destination buffer */

uLongf z_dst_nbytes = (uLongf)DEFLATE_SIZE_ADJUST(buf_size);
uLong z_src_nbytes = (uLong)buf_size;
int aggression = 9; /* Compression aggression setting */
uint32_t filter_mask = 0;
size_t buf_size = CHUNK_NX*CHUNK_NY*sizeof(int);

/* Create the data space */
if((dataspace = H5Screate_simple(RANK, dims, maxdims)) < 0)
    goto error;

/* Create a new file */
if((file = H5Fcreate(FILE_NAME5, H5F_ACC_TRUNC, H5P_DEFAULT, H5P_DEFAULT)) < 0)
    goto error;

/* Modify dataset creation properties, i.e. enable chunking and compression */
if((cparms = H5Pcreate(H5P_DATASET_CREATE)) < 0)
    goto error;

if((status = H5Pset_chunk(cparms, RANK, chunk_dims)) < 0)
    goto error;

if((status = H5Pset_deflate(cparms, aggression)) < 0)
    goto error;
```
goto error;

/* Create a new dataset within the file using cparms creation properties */
if((dset_id = H5Dcreate2(file, DATASETNAME, H5T_NATIVE_INT, dataspace, H5P_DEFAULT, cparms, H5P_DEFAULT)) < 0)
    goto error;

/* Initialize data for one chunk */
for(i = n = 0; i < CHUNK_NX; i++)
    for(j = 0; j < CHUNK_NY; j++)
        direct_buf[i][j] = n++;

/* Allocate output (compressed) buffer */
outbuf = malloc(z_dst_nbytes);
z_dst = (Bytef *)outbuf;

/* Perform compression from the source to the destination buffer */
ret = compress2(z_dst, &z_dst_nbytes, z_src, z_src_nbytes, aggression);

/* Check for various zlib errors */
if(Z_BUF_ERROR == ret) {
    fprintf(stderr, "overflow");
    goto error;
} else if(Z_MEM_ERROR == ret) {
    fprintf(stderr, "deflate memory error");
    goto error;
} else if(Z_OK != ret) {
    fprintf(stderr, "other deflate error");
    goto error;
}

/* Write the compressed chunk data repeatedly to cover all the chunks in the
 * dataset, using the direct write function. */
for(i=0; i<NX/CHUNK_NX; i++) {
    for(j=0; j<NY/CHUNK_NY; j++) {
        status = H5DOwrite_chunk(dset_id, dxpl_id, filter_mask, offset,
                                  z_dst_nbytes, outbuf);
offset[1] += CHUNK_NY;
}
offset[0] += CHUNK_NX;
offset[1] = 0;
}

/* Overwrite the first chunk with uncompressed data. Set the filter mask to 
* indicate the compression filter is skipped */
filter_mask = 0x00000001;
offset[0] = offset[1] = 0;
if(H5DOwrite_chunk(dset_id, dxpl_id, filter_mask, offset, buf_size, direct_buf) < 0)
    goto error;

/* Read the entire dataset back for data verification converting ints to longs*/
if(H5Dread(dataset, H5T_NATIVE_LONG, H5S_ALL, H5S_ALL, H5P_DEFAULT, 
         outbuf_long) < 0)
    goto error;

/* Data verification here */
:
:

...